

Quantum computers: the future attack that breaks today's messages

Daniel J. Bernstein & Tanja Lange

University of Illinois at Chicago & Technische Universiteit Eindhoven

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Cryptography

- ▶ Motivation #1: Communication channels are spying on our data.
- ▶ Motivation #2: Communication channels are modifying our data.



- ▶ Literal meaning of cryptography: "secret writing".
- ▶ Achieves various security goals by secretly transforming messages.

Cryptographic applications in daily life

- ▶ Mobile phones connecting to cell towers.
- ▶ Credit cards, EC-cards, access codes for banks.
- ▶ Electronic passports; electronic ID cards.
- ▶ Internet commerce, online tax declarations, webmail.
- ▶ Facebook, Gmail, WhatsApp, iMessage on iPhone.
- ▶ Any webpage with `https`.
- ▶ Encrypted file system on iPhone: see Apple vs. FBI.

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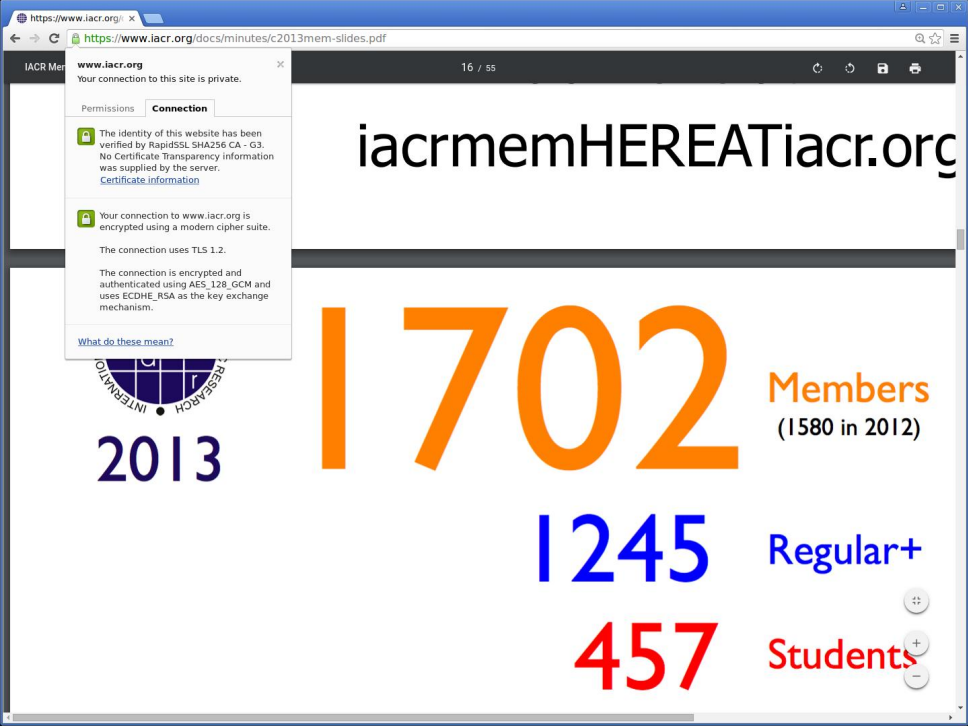
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Snowden in Reddit AmA

Arguing that you don't care about the right to privacy because you have nothing to hide is no different than saying you don't care about free speech because you have nothing to say.



www.iacr.org
Your connection to this site is private.

Permissions **Connection**

The identity of this website has been verified by RapidSSL SHA256 CA - G3. No Certificate Transparency information was supplied by the server.
[Certificate information](#)

Your connection to www.iacr.org is encrypted using a modern cipher suite.
The connection uses TLS 1.2.

The connection is encrypted and authenticated using AES_128_GCM and uses ECDHE_RSA as the key exchange mechanism.

[What do these mean?](#)

iacrmembership.org



1702

Members
(1580 in 2012)

1245

Regular+

457

Students

www.iacr.org



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iacrm



Cryptographic tools

Many factors influence the security and privacy of data:

- ▶ Secure storage, physical security; access control.
- ▶ Protection against alteration of data
⇒ [public-key signatures, message-authentication codes](#).
- ▶ Protection of sensitive content against reading
⇒ [encryption](#).

Many more security goals studied in cryptography

- ▶ Protecting against denial of service.
- ▶ Stopping traffic analysis.
- ▶ Securely tallying votes.
- ▶ Searching in and computing on encrypted data.
- ▶ ...

Cryptanalysis

- ▶ Cryptanalysis is the study of security of cryptosystems.
- ▶ Breaking a system can mean that the hardness assumption was not hard or that it just was not as hard as previously assumed.
- ▶ Public cryptanalysis is ultimately constructive – ensure that secure systems get used, not insecure ones.
- ▶ Weakened crypto ultimately backfires – attacks in 2018 because of crypto wars in the 90s.
- ▶ Good arsenal of general approaches to cryptanalysis. There are some automated tools.
- ▶ This area is constantly under development; researchers revisit systems continuously.





Security assumptions

- ▶ Hardness assumptions at the basis of all public-key and essentially all symmetric-key systems result from (failed) attempts at breaking systems. Security proofs are built only on top of those assumptions.
- ▶ A solid symmetric system is required to be as strong as exhaustive key search.
- ▶ For public-key systems the best attacks are faster than exhaustive key search. Parameters are chosen to ensure that the best attack is infeasible.

Key-size recommendations

	Parameter	Legacy	Future System Use	
			Near Term	Long Term
Symmetric Key Size	k	80	128	256
Hash Function Output Size	m	160	256	512
MAC Output Size*	m	80	128	256
RSA Problem	$\ell(n) \geq$	1024	3072	15360
Finite Field DLP	$\ell(p^n) \geq$	1024	3072	15360
	$\ell(p), \ell(q) \geq$	160	256	512
ECDLP	$\ell(q) \geq$	160	256	512
Pairing	$\ell(p^{k \cdot n}) \geq$	1024	6144	15360
	$\ell(p), \ell(q) \geq$	160	256	512

- ▶ Source: ECRYPT-CSA “[Algorithms, Key Size and Protocols Report](#)” (2018).
- ▶ These recommendations take into account attacks known today.
- ▶ Use extrapolations to larger problem sizes.
- ▶ Attacker power typically limited to 2^{128} operations (less for legacy).
- ▶ More to come on long-term security ...

Summary: current state of the art

- ▶ Currently used crypto (check the lock icon in your browser) starts with RSA, Diffie-Hellman (DH) in finite fields, or elliptic-curve Diffie-Hellman (ECDH).
- ▶ Older standards are RSA or elliptic curves from NIST (or Brainpool), e.g. NIST P256 or ECDSA.
- ▶ Internet currently moving over to [Curve25519](#) (Bernstein) and [Ed25519](#) (Bernstein, Duif, Lange, Schwabe, and Yang).
- ▶ For symmetric crypto TLS (the protocol behind https) uses AES or ChaCha20 and some MAC, e.g. AES-GCM or ChaCha20-Poly1305. High-end devices have support for AES-GCM, smaller ones do better with ChaCha20-Poly1305.
- ▶ Security is getting better. Some obstacles: bugs; untrustworthy hardware;

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- ▶ Security is getting better. Some obstacles: bugs; untrustworthy hardware; let alone anti-security measures such as AABill.



Algorithms for Quantum Computation: Discrete Logarithms and Factoring

Peter W. Shor
AT&T Bell Labs
Room 2D-149
600 Mountain Ave.
Murray Hill, NJ 07974, USA

Abstract

A computer is generally considered to be a universal computational device; i.e., it is believed able to simulate any physical computational device with a cost in computation time of at most a polynomial factor. It is not clear whether this is still true when quantum mechanics is taken into consideration. Several researchers, starting with David Deutsch, have developed models for quantum mechanical computers and have investigated their computational properties. This paper gives Las Vegas algorithms for finding discrete logarithms and factoring integers on a quantum computer that take a number of steps which is polynomial in the input size, e.g., the number of digits of the integer to be factored. These two problems are generally considered hard on a classical computer and have been used as the basis of several proposed cryptosystems. (We thus give the first examples of quantum cryptanalysis.)

[1, 2]. Although he did not ask whether quantum mechanics conferred extra power to computation, he did show that a Turing machine could be simulated by the reversible unitary evolution of a quantum process, which is a necessary prerequisite for quantum computation. Deutsch [9, 10] was the first to give an explicit model of quantum computation. He defined both quantum Turing machines and quantum circuits and investigated some of their properties.

The next part of this paper discusses how quantum computation relates to classical complexity classes. We will thus first give a brief intuitive discussion of complexity classes for those readers who do not have this background. There are generally two resources which limit the ability of computers to solve large problems: time and space (i.e., memory). The field of analysis of algorithms considers the asymptotic demands that algorithms make for these resources as a function of the problem size. Theoretical computer scientists generally classify algorithms as efficient when the number of steps of the algorithms grows as a polynomial in the size of the input. The class of prob-

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- ▶ Shor’s algorithm solves in polynomial time:
 - ▶ Integer factorization. RSA is dead.
 - ▶ The discrete-logarithm problem in finite fields. DSA is dead.
 - ▶ The discrete-logarithm problem on elliptic curves. ECDSA is dead.
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- ▶ This breaks all current public-key cryptography on the Internet!
- ▶ Also, Grover’s algorithm speeds up brute-force searches.
- ▶ Example: Only 2^{64} quantum operations to break AES-128;
 2^{128} quantum operations to break AES-256.

History of post-quantum cryptography

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- ▶ PQCrypto 2006: International Workshop on Post-Quantum Cryptography.

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- ▶ PQCrypto 2008, PQCrypto 2010, PQCrypto 2011, PQCrypto 2013.
- ▶ 2014 EU publishes H2020 call including post-quantum crypto as topic.
- ▶ ETSI working group on “Quantum-safe” crypto.
- ▶ PQCrypto 2014.
- ▶ April 2015 NIST hosts first workshop on post-quantum cryptography
- ▶ August 2015 NSA wakes up



NSA announcements

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Post-quantum becoming mainstream

- ▶ PQCrypto 2016: 22–26 Feb in Fukuoka, Japan, > 200 people



- ▶ NIST called for post-quantum proposals (deadline Nov 2017).
- ▶ 82 submissions; big effort to analyze, implement, prove, . . .



PQCrypto 2018
The Ninth International Conference on Post-Quantum Cryptography
Fort Lauderdale, Florida, April 9-11, 2018



Confidence-inspiring crypto takes time to build

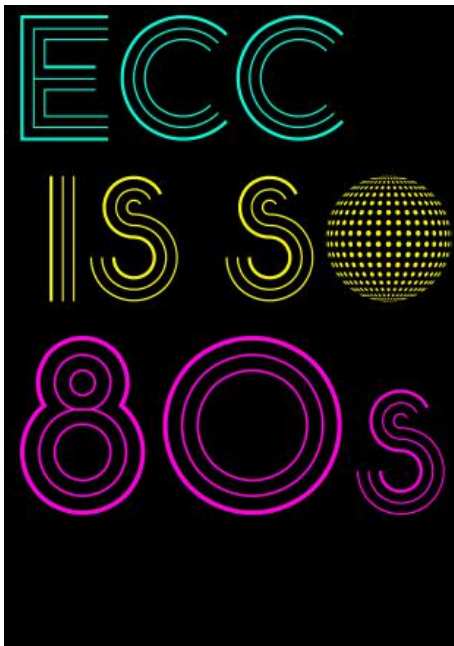
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 - ▶ Focus on implementations meeting performance requirements.
 - ▶ Integrate securely into real-world applications.

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 - ▶ Integrate securely into real-world applications.
- ▶ Example: ECC introduced **1985**; big advantages over RSA. Robust ECC started to take over the Internet in **2015**.
- ▶ Can't wait for quantum computers before finding a solution!



Even higher urgency for long-term confidentiality

- ▶ Today's encrypted communication is being stored by attackers and will be decrypted years later with quantum computers. Danger for human-rights workers, medical records, journalists, security research, legal proceedings, state secrets, . . .



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- ▶ Signature schemes can be replaced once a quantum computer is built – but there will not be a public announcement . . . and an important function of signatures is to protect operating system upgrades.
- ▶ Protect your upgrades *now* with post-quantum signatures.

Standardize now? Standardize later?

- ▶ Standardize now!
 - ▶ Rolling out crypto takes long time.
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- ▶ Standardize later!
 - ▶ Current options are not satisfactory.
 - ▶ Once rolled out, it's hard to change systems.
 - ▶ Please wait for the research results, will be much better!
- ▶ But what about users who rely on long-term secrecy of today's communication?
- ▶ Recommend now, standardize later. General roll out later.
- ▶ Recommend very conservative systems now; users who care will accept performance issues and gladly update to faster/smaller options later.
- ▶ But: Find out now where you rely on crypto; make an inventory.
- ▶ Important to raise awareness.

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- ▶ PQCRYPTO was an EU project in H2020, running 2015 – 2018.
- ▶ PQCRYPTO designed a portfolio of high-security post-quantum public-key systems, and improved the speed of these systems, adapting to the different performance challenges of mobile devices, the cloud, and the Internet.

Initial recommendations of long-term secure post-quantum systems

Daniel Augot, Lejla Batina, Daniel J. Bernstein, Joppe Bos,
Johannes Buchmann, Wouter Castryck, Orr Dunkelman,
Tim Güneysu, Shay Gueron, Andreas Hülsing,
Tanja Lange, Mohamed Saied Emam Mohamed,
Christian Rechberger, Peter Schwabe, Nicolas Sendrier,
Frederik Vercauteren, Bo-Yin Yang

Initial recommendations

- ▶ **Symmetric encryption** Thoroughly analyzed, 256-bit keys:
 - ▶ AES-256
 - ▶ Salsa20 with a 256-bit key

Evaluating: Serpent-256, ...

- ▶ **Symmetric authentication** Information-theoretic MACs:
 - ▶ GCM using a 96-bit nonce and a 128-bit authenticator
 - ▶ Poly1305

- ▶ **Public-key encryption** McEliece with binary Goppa codes:
 - ▶ length $n = 6960$, dimension $k = 5413$, $t = 119$ errors

Evaluating: QC-MDPC, Stehlé-Steinfeld NTRU, ...

- ▶ **Public-key signatures** Hash-based (minimal assumptions):
 - ▶ XMSS with any of the parameters specified in CFRG draft
 - ▶ SPHINCS-256

Evaluating: HFEv-, ...

Systems expected to survive

- ▶ Code-based crypto
- ▶ Hash-based signatures
- ▶ Isogeny-based crypto: new kid on the block, promising short keys and key exchange without communication (static-static) as possibility; needs more research on security; not covered here.
- ▶ Lattice-based crypto
- ▶ Multivariate crypto
- ▶ Symmetric crypto

Maybe some more, maybe some less.

Post-quantum secret-key authenticated encryption



- ▶ Very easy solutions if secret key k is long uniform random string:
 - ▶ “One-time pad” for encryption.
 - ▶ “Wegman–Carter MAC” for authentication.
- ▶ AES-256: Standardized method to expand 256-bit k into string indistinguishable from long k .
- ▶ AES introduced in 1998 by Daemen and Rijmen. Security analyzed in papers by dozens of cryptanalysts.
- ▶ No credible threat from quantum algorithms. Grover costs 2^{128} .
- ▶ Some recent results assume attacker has quantum access to computation, then some systems are weaker . . . but I’d know if my laptop had turned into a quantum computer.

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NIST Post-Quantum “Competition”

December 2016, after public feedback: NIST [calls for submissions](#) of post-quantum cryptosystems to standardize.

30 November 2017: NIST [receives 82 submissions](#).

Overview from Dustin Moody’s (NIST) talk at Asiacrypt 2017:

	Signatures	KEM/Encryption	Overall
Lattice-based	4	24	28
Code-based	5	19	24
Multi-variate	7	6	13
Hash-based	4		4
Other	3	10	13
Total	23	59	82

“Complete and proper” submissions

21 December 2017: NIST posts [69 submissions](#) from 260 people.

BIG QUAKE. BIKE. CFPKM. Classic McEliece. Compact LWE. CRYSTALS-DILITHIUM. CRYSTALS-KYBER. DAGS. Ding Key Exchange. DME. DRS. DualModeMS. Edon-K. EMBLEM and R.EMBLEM. FALCON. FrodoKEM. GeMSS. Giophantus. Gravity-SPHINCS. Guess Again. Gui. HILA5. HiMQ-3. HK17. HQC. KINDI. LAC. LAKE. LEDAkem. LEDApkc. Lepton. LIMA. Lizard. LOCKER. LOTUS. LUOV. McNie. Mersenne-756839. MQDSS. NewHope. NTRUencrypt. NTRU-HRSS-KEM. NTRU Prime. NTS-KEM. Odd Manhattan. OKCN/AKCN/CNKE. Ouroboros-R. Picnic. pqNTRUSign. pqRSA encryption. pqRSA signature. pqsigRM. QC-MDPC KEM. qTESLA. RaCoSS. Rainbow. Ramstake. RankSign. RLCE-KEM. Round2. RQC. RVB. SABER. SIKE. SPHINCS+. SRTPI. Three Bears. Titanium. WalnutDSA.

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BIG QUAKE. BIKE. CFPKM. Classic McEliece. Compact LWE. CRYSTALS-DILITHIUM. CRYSTALS-KYBER. DAGS. Ding Key Exchange. DME. DRS. DualModeMS. Edon-K. EMBLEM and R.EMBLEM. FALCON. FrodoKEM. GeMSS. Giophantus. Gravity-SPHINCS. Guess Again. Gui. HILA5. HiMQ-3. HK17. HQC. KINDI. LAC. LAKE. LEDAkem. LEDApkc. Lepton. LIMA. Lizard. LOCKER. LOTUS. LUOV. McNie. Mersenne-756839. MQDSS. NewHope. NTRUencrypt. NTRU-HRSS-KEM. NTRU Prime. NTS-KEM. Odd Manhattan. OKCN/AKCN/CNKE. Ouroboros-R. Picnic. pqNTRUSign. pqRSA encryption. pqRSA signature. pqsigRM. QC-MDPC KEM. qTESLA. RaCoSS. Rainbow. Ramstake. RankSign. RLCE-KEM. Round2. RQC. RVB. SABER. SIKE. SPHINCS+. SRTPI. Three Bears. Titanium. WalnutDSA.

Code-based encryption

BIG QUAKE

Classic McEliece

LAKE

LOCKER

DAGS

LEDAkem

LEDApkc

Lepton

McNie

Edon-K†

BIKE☢

HQC☢

NTS-KEM☢

Ouroboros-R☢

QC-MDPC KEM☢

RQC☢

RLCE-KEM☢

Code-based encryption

BIG QUAKE
Classic McEliece
LAKE
LOCKER
DAGS
LEDAkem
LEDApkc
Lepton
McNie

Edon-K†
BIKE☢
HQC☢
NTS-KEM☢
Ouroboros-R☢
QC-MDPC KEM☢
RQC☢
RLCE-KEM☢

Ouroboros-R☢, LAKE, LOCKER are merging into “ROLLO”.
LEDAkem and **LEDApkc** are also merging.

Code-based encryption

BIG QUAKE
Classic McEliece
LAKE
LOCKER
DAGS
LEDAkem
LEDApkc
Lepton
McNie

Edon-K†
BIKE‡
HQC‡
NTS-KEM‡
Ouroboros-R‡
QC-MDPC KEM‡
RQC‡
RLCE-KEM‡

Ouroboros-R‡, LAKE, LOCKER are merging into “ROLLO”.
LEDAkem and **LEDApkc** are also merging.

†: submitter has withdrawn submission.

‡: submitter has claimed patent on submission.

Warning: Other people could also claim patents.

Lattice-based encryption

CRYSTALS-KYBER
EMBLEM and R.EMBLEM
FrodoKEM
KINDI
LAC
LIMA
LOTUS
NewHope
NTRUEncrypt
NTRU-HRSS-KEM

NTRU Prime
Odd Manhattan
SABER
Titanium
HILA5
Ding Key Exchange🦊
Lizard🦊
KCL OKCN/AKCN/CNKE🦊
Round2🦊
Compact LWE🦊

HILA5, Round2🦊 are merging into “Round5” .
NTRUEncrypt, NTRU-HRSS-KEM are also merging.

Other encryption

SIKE: isogeny-based encryption

Other encryption

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Mersenne-756839: integer-ring encryption

Ramstake: integer-ring encryption

Three Bears: integer-ring encryption

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CFPKM: multivariate encryption

SRTPI†: multivariate encryption

DME☢: multivariate encryption

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Ramstake: integer-ring encryption

Three Bears: integer-ring encryption

pqRSA: factoring-based encryption

CFPKM: multivariate encryption

SRTPI†: multivariate encryption

DME♣: multivariate encryption

Guess Again: hard to classify

HK17†: hard to classify

RVB†: hard to classify

Signatures

Gravity-SPHINCS: hash-based

Picnic: hash-based

SPHINCS+: hash-based

DualModeMS: multivariate

GeMSS: multivariate

HiMQ-3: multivariate

LUOV: multivariate

Giophantus: multivariate

Gui[⚡]: multivariate

MQDSS[⚡]: multivariate

Rainbow[⚡]: multivariate

pqRSA: factoring-based

CRYSTALS-DILITHIUM: lattice-based

qTESLA: lattice-based

DRS: lattice-based

FALCON[⚡]: lattice-based

pqNTRUSign[⚡]: lattice-based

pqsigRM: code-based



RaCoSS: code-based


RankSign[†]: code-based

WalnutDSA[⚡]: braid-group

Post-quantum public-key signatures: hash-based



- ▶ Secret key , public key .
- ▶ Only one prerequisite: a good hash function, e.g. SHA3-512, ... Hash functions map long strings to fixed-length strings.

Signature schemes use hash functions in handling .

- ▶ Old idea: 1979 Lamport one-time signatures.
- ▶ 1979 Merkle extends to more signatures.

Pros and cons



Datatracker

Groups

Documents

Meetings

Other

User

Internet Research Task Force (IRTF)
Request for Comments: 8391
Category: Informational
ISSN: 2070-1721

A. Huelsing
TU Eindhoven
D. Butin
TU Darmstadt
S. Gazdag
genua GmbH
J. Rijneveld
Radboud University
A. Mohaisen
University of Central Florida
May 2018

XMSS: eXtended Merkle Signature Scheme

Pros:

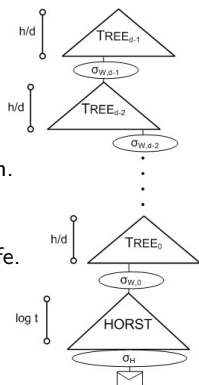
- ▶ Post quantum
- ▶ Only need secure hash function
- ▶ Small public key
- ▶ Security well understood
- ▶ Fast
- ▶ Accepted as [RFC 8391](#)

Cons:

- ▶ Biggish signature
- ▶ Stateful
Adam Langley “for most environments it’s a huge foot-cannon.”

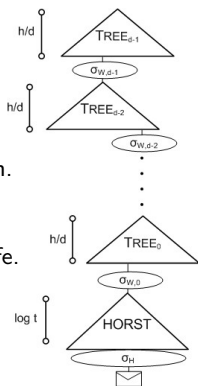
Stateless hash-based signatures

- ▶ Idea from 1987 Goldreich:
 - ▶ Signer builds huge tree of certificate authorities.
 - ▶ Signature includes certificate chain.
 - ▶ Each CA is a hash of master secret and tree position. This is deterministic, so don't need to store results.
 - ▶ **Random** bottom-level CA signs message. Many bottom-level CAs, so one-time signature is safe.



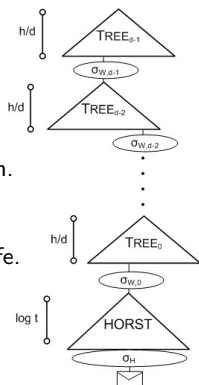
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- ▶ 1.2 MB: average Debian package size.
- ▶ 1.8 MB: average web page in Alexa Top 1000000.
- ▶ 0.041 MB: SPHINCS signature, new optimization of Goldreich. Modular, guaranteed as strong as its components (hash, PRNG). Well-known components chosen for 2^{128} post-quantum security.
sphincs.cr.yp.to

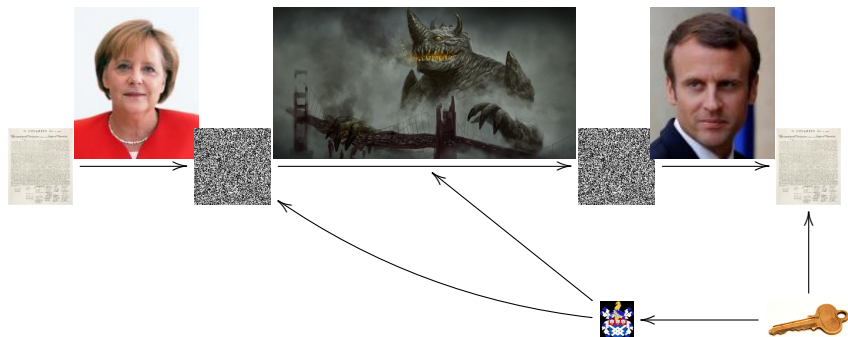




NIST submission SPHINCS+

- ▶ Same as SPHINCS in terms of high level scheme design, but better few-time signatures.
- ▶ New protection against multi-target attacks.
- ▶ New few-time signature scheme FORS instead of HORST (different way of combining Merkle trees).
- ▶ Smaller signatures – 30kB instead of 41kB – or more signatures.
- ▶ Smaller public keys.
- ▶ Three versions (different hash functions)
 - ▶ SPHINCS+-SHA3 (using SHAKE256),
 - ▶ SPHINCS+-SHA2 (using SHA-256),
 - ▶ SPHINCS+-Haraka (using the Haraka short-input hash function).

See <https://sphincs.org/> for more details.

Post-quantum public-key encryption: code-based



- ▶ Alice uses Bob's public key  to encrypt.
- ▶ Bob uses his secret key  to decrypt.
- ▶ Code-based crypto proposed by McEliece in 1978 using Goppa codes.
- ▶ Almost as old as RSA, but much stronger security history.
- ▶ Many further improvements, e.g. Niederreiter system for smaller keys.

One-wayness (OW-CPA)

Fundamental security question:

Given random parity-check matrix H and syndrome s ,
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guiding sizes in 1978 McEliece.

The McEliece system (with later key-size optimizations)
uses $(c_0 + o(1))\lambda^2(\lg \lambda)^2$ -bit keys as $\lambda \rightarrow \infty$
to achieve 2^λ security against Prange's attack.

Here $c_0 \approx 0.7418860694$.

40 years and more than 30 analysis papers later

1962 Prange; 1981 Clark–Cain, crediting Omura; 1988 Lee–Brickell; 1988 Leon; 1989 Krouk; 1989 Stern; 1989 Dumer; 1990 Coffey–Goodman; 1990 van Tilburg; 1991 Dumer; 1991 Coffey–Goodman–Farrell; 1993 Chabanne–Courteau; 1993 Chabaud; 1994 van Tilburg; 1994 Canteaut–Chabanne; 1998 Canteaut–Chabaud; 1998 Canteaut–Sendrier; 2008 Bernstein–Lange–Peters; 2009 Bernstein–Lange–Peters–van Tilborg; 2009 Bernstein (**post-quantum**); 2009 Finiasz–Sendrier; 2010 Bernstein–Lange–Peters; 2011 May–Meurer–Thomae; 2012 Becker–Joux–May–Meurer; 2013 Hamdaoui–Sendrier; 2015 May–Ozerov; 2016 Canto Torres–Sendrier; 2017 Kachigar–Tillich (**post-quantum**); 2017 Both–May; 2018 Both–May; 2018 Kirshanova (**post-quantum**).

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Same $c_0 \approx 0.7418860694$.

Replacing λ with 2λ stops all known *quantum* attacks.

NIST submission Classic McEliece

- ▶ Security asymptotics unchanged by 40 years of cryptanalysis.
- ▶ Short ciphertexts.
- ▶ Efficient and straightforward conversion of OW-CPA PKE into IND-CCA2 KEM.
- ▶ Constant-time software implementations.
- ▶ FPGA implementation of full cryptosystem.
- ▶ Open-source (public domain) implementations.
- ▶ No patents.

Metric	mceliece6960119	mceliece8192128
Public-key size	1047319 bytes	1357824 bytes
Secret-key size	13908 bytes	14080 bytes
Ciphertext size	226 bytes	240 bytes
Key-generation time	1108833108 cycles	1173074192 cycles
Encapsulation time	153940 cycles	188520 cycles
Decapsulation time	318088 cycles	343756 cycles

See <https://classic.mceliece.org> for more details.

NIST submission NTRU Prime

- ▶ Lattice-based encryption – smaller public keys.
- ▶ Less structure for the attacker to use:
 - ▶ Computation is done modulo prime instead of modulo power of 2.
 - ▶ Rings change from using polynomial $x^n - 1$ or $x^n + 1$ to $x^p - x - 1$, p prime.
 - ▶ No (nontrivial) subrings or fields.
- ▶ No decryption failures.

Metric	Streamlined NTRU Prime 4591⁷⁶¹	NTRU LPrime 4591⁷⁶¹
Public-key size	1218 bytes	1047 bytes
Secret-key size	1600 bytes	1238 bytes
Ciphertext size	1047 bytes	1175 bytes
Key-generation time	5925834 cycles	44940 cycles
Encapsulation time	45468 cycles	80596 cycles
Decapsulation time	94744 cycles	113272 cycles

See <https://ntruprime.cr.yp.to/> for more details.

Links and upcoming events

- ▶ <https://csrc.nist.gov/projects/post-quantum-cryptography/round-1-submissions>: NIST PQC competition.
- ▶ Early January 2019: NIST announces second-round candidates.
- ▶ 1 & 2 July 2019: Executive summer school in Eindhoven.
- ▶ <https://pqcrypto.eu.org>: PQCRYPTO EU project.
 - ▶ Expert [recommendations](#).
 - ▶ Free software libraries ([libpqcrypto](#), [pqm4](#), [pqhw](#)).
 - ▶ Lots of reports, scientific papers, (overview) presentations.
- ▶ <https://2017.pqcrypto.org/school>: PQCRYPTO summer school with 21 lectures on video + slides + exercises.
- ▶ <https://2017.pqcrypto.org/exec>: Executive school (12 lectures), less math, more overview. So far slides, soon videos.
- ▶ [PQCrypto 2017](#) conference.
- ▶ [PQCrypto 2016](#) with slides and videos from lectures + school.
- ▶ <https://pqcrypto.org>: Our survey site.
 - ▶ Many pointers: e.g., PQCrypto conference series.
 - ▶ Bibliography for 4 major PQC systems.